FEC Decoding

## Keywords

- FEC
- Trellis
- Viterbi
- CC1100
- CC1100E
- CC1101
- CC1110Fx
- CC1111Fx
- CC1150
- CC2500
- CC2510Fx
- CC2511Fx
- CC2550
- CC430FX


## 1 Introduction

The CC1100, CC1100E, CC1101, CC1110Fx, CC1111Fx, CC1150, CC2500, CC2510Fx, CC2511Fx, and CC2550 all implement FEC encoding and decoding in HW. The purpose of this design note is to describe how one can implement the same FEC decoding in SW. This is in particular very important for the CC430Fx
device, which has the same radio as the CC1101 and CC1110/11Fx but without HW FEC included. This design note is not meant as a tutorial on FEC and it will not cover the FEC encoding, as it is described in
DN504 [1].

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2 Abbreviations

| CRC | Cyclic Redundancy Check |
| :--- | :--- |
| FEC | Forward Error Correction |
| FIFO | First In First Out |
| HW | Hardware |
| LSB | Least Significant Bit |
| MSB | Most Significant Bit |
| RAM | Random Access Memory |
| SW | Software |

3 Implementation

### 3.1 Code Example Assumptions and Limitations

Assume that you want to use the CC1101 [5] to transmit a packet and the CC430Fx [2] to receive it. The payload is 29 bytes and 2 bytes of CRC are appended. If FEC is enabled (MDMCFG1.FEC_EN = 1) on the transmitter, 64 bytes will be transmitted over the air in addition to preamble and sync word. Due to the appended trellis terminator and the size of the interleaving buffer the packet length will always be a multiple of 4 after encoding (see DN504 [1]). The number of bytes sent over the air (not including preamble and sync word) can be calculated as shown in Equation 1.

$$
\text { \# of Bytes on the Air }=\left(\left(\frac{\text { Payload Length }+2 \text { Bytes Optional CRC }}{2}\right)+1\right) \cdot 4
$$

## Equation 1. \# of Bytes on the Air ${ }^{1}$

The code example shown in Section 3.2 does not show how to set up the CC430Fx to receive a packet and it does not show how to implement the function (readRxFifo) that will read from the RXFIFO in Figure 5 (see the CC430Fx User's Guide [2] on how this can be done). It is assumed that a flag, packetReceived, is asserted when the packet is received and the 64 bytes are in the RXFIFO. It is important to notice that there is no packet size limitation to the FEC decoding itself, but if more than 64 bytes are sent over the air, the receiver must start to read the RXFIFO before the complete packet is received and additional RAM is needed to store the un-coded packet. When 64 bytes or less is sent on the air, the un-coded data can simply be kept in the RXFIFO until being decoded. The CC1101 [5] only support fixed packet length mode (PKTCTRLO. LENGTH_CONFIG = 0) when FEC is enabled, so the CC430Fx should also use this mode. Overflow of the RXFIFO will therefore not be an issue as long as the maximum packet length is less than 64 bytes (PKTLEN <= $0 \times 40$ ).

If the receiver is not the CC430Fx but some other radio which do not have an RXFIFO, it will be necessary to store the receive packet in a temporary buffer and then the function readRxFifo should read from this buffer instead of from the RXFIFO. It is assumed that the bytes in this buffer are read in the same order as they would have been read from a traditional FIFO.

[^0]
### 3.2 C Code

The C code in this section is organized as follows: Figure 1; Function Prototypes, Global Variables, and Defines, Figure 2; Function Definitions, Figure 3 and Figure 4; FEC Decoder Implementation, and Figure 5; main.

```
* FUNCTION PROTOTYPES
unsigned short fecDecode(unsigned char *pDecData, unsigned char* pInData, unsigned short RemBytes);
static unsigned char hammWeight(unsigned char a);
static unsigned char min(unsigned char a, unsigned char b);
static unsigned short calcCRC(unsigned char crcData, unsigned short crcReg);
GLOBAL VARIABLES
* GI
// The payload + CRC are 31 bytes. This way the complete packet to be received will fit in the RXFIFO
unsigned char rxBuffer[4]; // Buffer used to hold data read from the RXFIFO (4 bytes are read at a time)
unsigned char rxBuffer[4];
// Look-up source state index when
// Destination state --\ /-- Each of two possible source states
const unsigned char aTrellisSourceStateLut[8] [2] =
\begin{tabular}{|c|c|c|c|}
\hline \(\{0,4\}\), & // State & \{0,4\} & -> State \\
\hline \(\{0,4\}\), & // State & 0,4\} & State \\
\hline \(\{1,5\}\), & // State & \{1,5\} & State \\
\hline \(\{1,5\}\), & // State & \{1,5\} & -> State \\
\hline \(\{2,6\}\), & // State & \(\{2,6\}\) & -> State \\
\hline \(\{2,6\}\), & // State & \{2, 6\} & State \\
\hline \{3, 7\}, & // State & \{3,7\} & State \\
\hline \{3, 7\}, & // State & \{3, 7 \} & -> State \\
\hline
\end{tabular}
};
7/ Look-up expected output when:
Destination state --\ /-- Each of two possible source states
const unsigned char aTrellisTransitionOutput[8][2] =
|
\begin{tabular}{|c|c|c|c|c|c|c|}
\hline \(0,3\}\), & // State & \{0, 4 \} & te & es & \{"00" & \\
\hline \(3,0\}\) & // State & \{0, 4 \} & -> State 1 & produces & \{"11" & 00" \(\}\) \\
\hline 1, 2\(\}\) & // State & \{1,5\} & State 2 & produce & "01" & 10 \\
\hline 2, 1\} & // State & \{1,5\} & State 3 & produces & "10 & "01" \\
\hline \(\{3,0\}\) & // State & \{2, 6\} & State 4 & produces & "11" & 00 \\
\hline 0,3 & // State & 2,6\} & State 5 & produces & "00" & 11 \\
\hline \(2,1\}\) & // State & 3, 7 \(\}\) & -> State 6 & produces & "10" & 01 \\
\hline , 2\} & // State & \{3, 7\} & -> State 7 & produces & "01" & 10 \\
\hline
\end{tabular}
};
// Look-up input bit at encoder when
// Destination state --
const unsigned char aTrellisTransitionInput[8] =
cons
0,
1,
1,
1,
};
/*********
*/
// NUMBER_OF_BYTES_AFTER_DECODING should be given the length of the payload + CRC (CRC is optional)
#define NUMMEER OF BYTES AFTER DECODING 31
#define NUMBER_OF_BYTES_AFTER_DECODING 31 ( (NUMBER_OF_BYTES_AFTER_DECODING / 2) + 1))
```

Figure 1. Function Prototypes, Global Variables, and Defines

```
l************************************************************************************************************************
* @brief Calculates Hamming weight of byte (# bits set)
* @param a - Byte to find the Hamming weight for
* @return Hamming weight (# of bits set in a)
static unsigned char hammWeight (unsigned char a)
    a=((a&0xAA) >> 1) + (a& & 0\times55);
    a = ((a& & xCC) >> 2) + (a& & 0x33)
    =((a& & 0xF0) >> 4) + (a & 0x0F)
    return a
}
/******************
@brief Returns the minimum of two values
@param a - Value 1
@return Minimum of two values
    @return (Vinimum of two values 
* Value 2 (Value 2 < Value 1)
static unsigned char min(unsigned char a, unsigned char b)
    return (a <= b ? a : b);
}
@fn calcCRC
@brief Calculates a checksum over n data bytes
    Example of usage
    checksum = 0xFFFF;
    checksum = calcCRC(dataBytes[i], checksum)
@param crcData - checksum (initially set to 0xFFFF)
    crcReg - data byte
*/@return Checksum
static unsigned short calcCRC(unsigned char crcData, unsigned short crcReg)
    unsigned char 1;
    for (i = 0; i < 8; i++) {
        if (((crcReg & 0x8000) >> 8) ^ (crcData & 0x80))
        crcReg = (crcReg << 1) ^ 0x8005;
        else
        crcReg = (crcReg << 1);
        crcData <<= 1;
    }
    return crcReg;
}
```

Figure 2. Function Definitions

```
l************************************************************************************************************************
@brief De-interleaves and decodes a given input buffer
@ @param pDecData - Pointer to where to put decoded data (NULL when initializing at start of packet)
    Pointer to received data
    nRemBytes - of remaining (decoded) bytes to decode
@return Number of bytes of decoded data stored at pDecData
unsigned short fecDecode(unsigned char *pDecData, unsigned char* pInData, unsigned short nRemBytes)
// Two sets of buffers (last, current) for each destination state for holding:
    static unsigned char nCost[2] [8]; // Accumulated path cost
```



```
    // Indices of (last, current) buffer for each iteration
    tatic unsigned char iLastBuf
    static unsigned char iCurrBuf;
    // Number of bits in each path buffer
    static unsigned char nPathBits;
    // Variables used to hold # Viterbi iterations to run, # bytes output,
    // minimum cost for any destination state, bit index of input symbol
    unsigned char nIterations;
    unsigned short nOutputBytes = 0;
    insigned char nuinCost;
    // Initialize variables at start of packet (and return without doing any more)
    f (pDecData == NULL) {
        unsigned char n ;
        memset (nCost, 0, sizeof(nCost));
        for (n = 1; n < 8; n++)
            nCost[0][n] = 100;
        iLastBuf = 0
        iCurrBuf = 1;
        nPathBits = 0;
        return 0;
    }
    {
        unsigned char aDeintData[4];
        signed char iOut
        signed char iIn;
        // De-interleave received data (and change pInData to point to de-interleaved data)
        for (iOut = 0; iOut < 4; iOut++) {
            unsigned char dataByte = 0;
            for (iIn = 3; iIn >= 0; iIn--)
            dataByte = (dataByte << 2) | ((pInData[iIn] >>( 2 * iOut)) & 0x03);
            aDeintData[iOut] = dataByte;
    }
    pInData = aDeintData;
}
// Process up to 4 bytes of de-interleaved input data, processing one encoder symbol (2b) at a time
for (nIterations = 16; nIterations > 0; nIterations--) {
    unsigned char iDestState;
    unsigned char symbol = ((*pInData) >> iBit) & 0x03;
    // Find minimum cost so that we can normalize costs (only last iteration used)
    nMinCost = 0xFF;
    // Get 2b input symbol (MSB first) and do one iteration of Viterbi decoding
    if ((iBit -= 2) < 0) {
        iBit = 6; ./| Update pointer to the next byte of received data
        pInData++; // Update pointer to the next byte of received data
    }
    // For each destination state in the trellis, calculate hamming costs for both possible paths into state and
    / select the one with lowest cost
    for (iDestState = 0; iDestState < 8; iDestState++) {
        unsigned char ncost0;
        unsigned char nCost1;
        unsigned char iSrcState0;
        unsigned char iSrcStatel;
        unsigned char nInputBit;
        nInputBit = aTrellisTransitionInput[iDestState];
        // Calculate cost of transition from each of the two source states (cost is Hamming difference between
        // received 2b symbol and expected symbol for transition)
    iSrcState0 = aTrellisSourceStateLut[iDestState] [0]
    Cost0 = nCost[iLastBuf][iSrcState0]
    nCost0 += hammWeight(symbol ^ aTrellisTransitionOutput[iDestState][0])
    iSrcState1 = aTrellisSourceStateLut[iDestState][1]
    nCost1 = nCost[iLastBuf][iSrcState1]
    nCost1 += hammWeight(symbol ^ aTrellisTransitionOutput[iDestState][1])
```

Figure 3. FEC Decoder Implementation (1)

```
    // select transition that gives lowest cost in destination state, copy that source state's path and add
    // new decoded bit
    if (nCost0 <= nCost1) {
        nCost[iCurrBuf][iDestState] = nCost0;
        nMinCost = min(nMinCost, nCost0);
        aPath[iCurrBuf][iDestState] = (aPath[iLastBuf][iSrcState0] << 1) | nInputBit;
    } else
        nCost[iCurrBuf][iDestState] = nCost1;
        nMinCost = min(nMinCost, nCost1);
        aPath[iCurrBuf][iDestState] = (aPath[iLastBuf][iSrcState1] << 1) | nInputBit;
    }
    }
    nPathBits++;
    // If trellis history is sufficiently long, output a byte of decoded data
    if (nPathBits == 32) {
        *pDecData++ = (aPath[iCurrBuf][0] >> 24) & OxFF;
        nOutputBytes++;
        nPathBits --
    }
    // After having processed 3-symbol trellis terminator, flush out remaining data
    // After having processed 3-symbol trellis terminator, flush out
    if ((nRemBytes <= 3) && (nP
        *pDecData++ = (aPath[iCurrBuf][0] >> (nPathBits - 8)) & 0xFF;
            nOutputBytes++;
            nPathBits -= 8;
        }
    return nOutputBytes;
    }
    // Swap current and last buffers for next iteration
    iLastBuf = (iLastBuf + 1) % 2;
    iCurrBuf = (iCurrBuf + 1) % 2;
}
// Normalize costs so that minimum cost becomes 0
{
    unsigned char iState;
    for (iState = 0; iState < 8; iState++)
    nCost[iLastBuf][iState] -= nMinCost;
}
return nOutputBytes;
```

$\}^{r}$

Figure 4. FEC Decoder Implementation (2)

```
*)
    @ @fn main
    @brief This code example demonstrates how the fecDecode function can be used. It is assumed that a
    flag, packetReceived, is asserted when a packet is received (there are 64 bytes in the RXFIFO)
    * @param None
    @return None
void main(void)
    unsigned short checksum;
    n
    unsigned short nBytes; // Destination for decoded data
    // Init MCU and Radio
    while (1) {
    while (!packetReceived); // Wait for packet to be received (64 bytes in the RXFIFO)
    packetReceived = 0;
    pDecData = rxPacket;
    // Perform de-interleaving and decoding (both done in the same function)
    fecDecode(NULL, NULL, 0); // The function needs to be called with a NULL pointer for
    nBytes = NUMBER_OF_BYTES_AFTER_DECODING;
    while (nBytes >- 0)- {
        unsigned short nBytesOut;
        readRxFifo(RF_RXFIFO, rxBuffer, 4); // Read 4 bytes from the RXFIFO and store them in rxBuffer
        nBytesOut = fececDecode(pDecData, rxBuffer, nBytes);
        nBytes -= nBytesOut;
        pDecData += nBytesOut;
    }
    // Perform CRC check (Optional)
        unsigned short i;
        nBytes = NUMBER_OF_BYTES_AFTER_DECODING;
        checksum = 0xFFFF;
        for (i = 0; i < nBytes; i++
            checksum = calcCRC(rxPacket[i], checksum)
        if (!checksum) {
            // Do something to indicate that the CRC is OK
        }
    }
}
```

Figure 5. main

## 4 Explanation to the Code

The most important part of the code is the decoder part implemented in the function fecDecode (see Figure 3 and Figure 4). The function will process 4 and 4 bytes of received data since this is the data size the interleaver works on. This means that in most cases the function will be called several times for each received packet. The pseudo code for the function is shown in Figure 6.

```
fecDecode()
    / Variable Declaration
    // Initialize variables at start of packet (and return without doing any more)
    // De-interleave 4 bytes of received data (4 bytes of data means 16 (2b) encode symbols
    // For all 16 symbols do one iteration of Viterbi decoding
    for (nIterations = 16; nIterations > 0; nIterations--) {
        // Get 2b input symbol (MSB first) and do one iteration of Viterbi decoding
    // For each destination state in the trellis
    for (iDestState = 0; iDestState < 8; iDestState++) {
        // Calculate cost of transition from each of the two source states (cost is Hamming difference between
        // received 2b symbol and expected symbol for transition)
        // Select transition that gives lowest cost in destination state, copy that source state's path and add
        // new decoded bit
    }
    nPathBits++;
    // If trellis history is sufficiently long, output a byte of decoded data. After 32 iterations
    // (nPathBits == 32) the 8 MSB bits will be the same for all }8\mathrm{ surviving paths and a byte can be output
    // After having processed 3-symbol trellis terminator, flush out remaining data
    // Swap current and last buffers for next iteration
    }
    / Normalize costs so that minimum cost becomes 0
}
```

Figure 6. Pseudo Code for the FEC Encoder
The "key" elements of the code are the two for loops that for every symbol goes through each destination state in the trellis. An example is used to illustrate what is going on in this loop.

Example: A packet consisting of 5 bytes ( $0 \times 01,0 \times 02,0 \times 03,0 \times 04,0 \times 05$ ) is being interleaved and encoded by the CC1101 (MDMCFG1.FEC_EN = 1). The data transmitted on the air will be the following (preamble and sync word is not shown):
$0 \times 4 \mathrm{C}, 0 \times \mathrm{F} 0,0 \times 30,0 \times 10,0 \times \mathrm{C} 8,0 \times 7 \mathrm{C}, 0 \times \mathrm{C} 3,0 \times 23,0 \times 40,0 \times 34,0 \times 7 \mathrm{C}, 0 \times E 3$ (see DN504 [1])
In chunks of 4 and 4 bytes, this data will on the receiver side be interleaved, giving the symbols shown in Table 1 to be decoded (only the 4 first bytes are shown):

| Symbol \# | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 | 12 | 13 | 14 | 15 | 16 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| Symbol | 00 | 00 | 00 | 00 | 00 | 00 | 00 | 11 | 01 | 11 | 11 | 00 | 00 | 00 | 11 | 01 |
| Byte | 0x00 |  |  |  | $0 \times 03$ |  |  |  | 0x7C |  |  |  | 0x0D |  |  |  |

Table 1. Symbols to be Encoded


Figure 7. Trellis 1 (the number in () indicates the source state of that given cost)
For each received symbol (2b), all possible encoder output symbols (00, 01, 10, and 11) in Figure 7 are compared against the received symbol and a transition cost is calculated (nCost0 and nCost1). The appropriate transition cost is added to the accumulated path cost of each path that terminates in the source state on the left in the figure. It can be seen that there are two transitions into each destination state on the right in the figure. For each destination state the incoming transition with the lowest accumulated path cost is selected (the survivor path) and the other one thrown away - nothing is lost as all future paths that go through this state at this point in the trellis would do the same selection. Thus the number of paths that the Viterbi Algorithm tracks is always constant and the optimal path is always one of them.

00


Figure 8. Trellis 2

```
iCurrBuf = 1
iLastBuf = 0
nCost[iCurrBuf][] = [0, 2, 101, 101, 100, 100, 101, 101]
aPath[iCurrBuf][iDeststate] = (aPath[iLastBuf][iSrcState] << 1) | nInputBit;
aPath[1][0] = (aPath[0][0] << 1) Input to state 0 = 000000000000000000000000000000000000
aPath[1][1] = (aPath[0][0] << 1) Input to state 1 = 00000000000000000000000000000001
aPath[1][2] = (aPath[0][1] << 1) Input to state 2 = 00000000000000000000000000000000
aPath[1][3] = (aPath[0][1] << 1) Input to state 3 = 0000000000000000000000000000000001
aPath[1][4] = (aPath[0][6] << 1) Input to state 4 = 000000000000000000000000000000000000
aPath[1][5] = (aPath[0][2] << 1) Input to state 5 = 00000000000000000000000000000001
aPath[1][6] = (aPath[0][3] << 1) Input to state 6=00000000000000000000000000000000000
aPath[1][7] = (aPath[0][3] << 1) Input to state 7 = 00000000000000000000000000000001
```

Figure 9. $n$ Cost and aPath after $1^{\text {st }}$ Iteration (received symbol: $\mathbf{0 0}_{b}$ )

00
00


Figure 10. Trellis 3

```
iCurrBuf = 0
iLastBuf = 1
nCost[iCurrBuf] [] = [0, 2, 3, 3, 101, 101, 102, 102]
aPath[iCurrBuf] [iDestState] = (aPath[iLastBuf][iSrcState] << 1) | nInputBit;
aPath[0][0] = (aPath[1][0] << 1) Input to state 0 = 00000000000000000000000000000000000
aPath[0][1] = (aPath[1][0] << 1) Input to state 1 = 00000000000000000000000000000001
aPath[0][2] = (aPath[1][1] << 1) Input to state 2 = 00000000000000000000000000000010
aPath[0][3] = (aPath[1][1] << 1) Input to state 3 = 0000000000000000000000000000000011
aPath[0][4] = (aPath[1][6] << 1) Input to state 4 = 00000000000000000000000000000000
aPath[0][5] = (aPath[1][2] << 1) Input to state 5 = 00000000000000000000000000000001
aPath[0][6] = (aPath[1][3] << 1) Input to state 6=000000000000000000000000000000010
aPath[0][7] = (aPath[1][3] << 1) Input to state 7 = 00000000000000000000000000000011
```

Figure 11. nCost and aPath after $\mathbf{2}^{\text {nd }}$ Iteration (received symbol: $\mathbf{0 0}_{\mathrm{b}}$ )

00


Figure 12. Trellis 4

```
iCurrBuf = 1
iLastBuf = 0
nCost[iCurrBuf][] = [0, 2, 3, 3, 5, 3, 4, 4]
aPath[iCurrBuf][iDeststate] = (aPath[iLastBuf][iSrcState] << 1) | nInputBit;
aPath[1][0] = (aPath[0][0] << 1) | Input to state 0 = 00000000000000000000000000000000000
aPath[1][1] = (aPath[0][0] << 1) Input to state 1 = 00000000000000000000000000000001
aPath[1][2] = (aPath[0][1] << 1) Input to state 2 = 00000000000000000000000000000010
aPath[1][3] = (aPath[0][1] << 1) Input to state 3=000000000000000000000000000000011
aPath[1][4] = (aPath[0][2] << 1) Input to state 4 = 00000000000000000000000000000100
aPath[1][5] = (aPath[0][2] << 1) Input to state 5 = 00000000000000000000000000000101
aPath[1][6] = (aPath[0][3] << 1) Input to state 6=000000000000000000000000000000110
aPath[1][7] = (aPath[0][3] << 1) Input to state 7 = 00000000000000000000000000000111
```

Figure 13. nCost and aPath after 3rd Iteration (received symbol: 00 ${ }_{b}$ )


Figure 14. Trellis 5

```
iCurrBuf = 0
iLastBuf = 1
nCost[iCurrBuf][] = [0, 2, 3, 3, 4, 3, 4, 4]
aPath[iCurrBuf][iDeststate] = (aPath[iLastBuf][iSrcState] << 1) | nInputBit;
aPath[0][0] = (aPath[1][0] << 1) | Input to state 0 = 00000000000000000000000000000000000
aPath[0][1] = (aPath[1][0] << 1) Input to state 1 = 00000000000000000000000000000001
aPath[0][2] = (aPath[1][1] << 1) Input to state 2 = 00000000000000000000000000000010
aPath[0][3] = (aPath[1][1] << 1) Input to state 3 = 0000000000000000000000000000000011
aPath[0][4] = (aPath[1][6] << 1) Input to state 4 = 00000000000000000000000000001100
aPath[0][5] = (aPath[1][2] << 1) Input to state 5 = 00000000000000000000000000000101
aPath[0][6] = (aPath[1][3] << 1) Input to state 6=000000000000000000000000000000110
aPath[0][7] = (aPath[1][3] << 1) Input to state 7 = 00000000000000000000000000000111
```

Figure 15. $n$ Cost and aPath after $4^{\text {th }}$ Iteration (received symbol: $\mathbf{0 0}{ }_{b}$ )


Figure 16. Trellis 6

```
iCurrBuf = 1
iLastBuf = 0
nCost[iCurrBuf][] = [0, 2, 3, 3, 4, 3, 4, 4]
aPath[iCurrBuf][iDestState] = (aPath[iLastBuf][iSrcState] << 1) | nInputBit;
aPath[1][0] = aPath[0][0] << 1 Input to state 0 = 000000000000000000000000000000000000
aPath[1][1] = aPath[0][0] << 1 Input to state 1 = 00000000000000000000000000000001
aPath[1][2] = aPath[0][1] << 1 Input to state 2 = 00000000000000000000000000000010
aPath[1][3] = aPath[0][1] << 1 Input to state 3 = 0000000000000000000000000000000011
aPath[1][4] = aPath[0][6] << 1 Input to state 4 = 00000000000000000000000000001100
aPath[1][5] = aPath[0][2] << 1 Input to state 5 = 00000000000000000000000000000101
aPath[1][6] = aPath[0][3] << 1 Input to state 6=00000000000000000000000000000000110
aPath[1][7] = aPath[0][3] << 1 Input to state 7 = 00000000000000000000000000000111
```

Figure 17. nCost and aPath after $5^{\text {st }}$ Iteration (received symbol: $\mathbf{0 0}{ }_{b}$ )

After having processed 32 symbol, the 8 MSBs of $\mathrm{aPath}[0][0]\left(0000001_{\mathrm{b}}\right)$ is copied to rxFifo[0].

```
aPath[iCurrBuf][iDestState] = (aPath[iLastBuf][iSrcState] << 1) | nInputBit;
aPath[0][0] = 00000001000000100000001100001000
aPath[0][1] = 00000001000000100000001100000001
aPath[0][2] = 00000001000000100000001100000010
aPath[0][3] = 00000001000000100000001100000011
aPath[0][4] = 000000010000001000000001100000100
aPath[0][5] = 00000001000000100000001100000101
aPath[0][6] = 00000001000000100000001100000110
aPath[0][7] = 00000001000000100000001100000111
```

Figure 18. aPath after having Processed 32 Symbols
After 8 more symbols, aPath[iCurrBuf][iDestState] looks like in Figure 19 and 00000010b (8 MSBs of aPath[0][0]) are copied to rxFifo[1].

```
aPath[iCurrBuf][iDestState] = (aPath[iLastBuf][iSrcState] << 1) | nInputBit;
aPath[0][0] = 000000100000001100000100000000000
aPath[0][1] = 000000100000001100000100000001001
aPath[0][2] = 00000010000000110000010000000010
aPath[0][3] = 00000010000000110000010000000011
aPath[0][4] = 000000100000001100000100000000100
aPath[0][5] = 00000010000000110000010000000101
aPath[0][6] = 000000100000001100000100000000110
aPath[0][7] = 000000100000001100000100000000111
```

Figure 19. aPath after having Processed 40 Symbols
After having processed 3 more symbols, the trellis terminator has been processed and the remaining bytes of the packet are being copied to rxFifo.
rxFifo[2] = aPath[1][0] bit 26:19 $=00000011_{\text {b }}$
rxFifo[3] = aPath[1][0] bit 18:11 $=00000100_{b}$
rxFifo[4] $=\mathbf{a P a t h [ 1 ] [ 0 ] ~ b i t ~ 1 0 : 3 ~}=00000101_{b}$

```
aPath[iCurrBuf][iDestState] = (aPath[iLastBuf][iSrcState] << 1) nInputBit;
aPath[1][0] = 000100000001100000100000000101000
aPath[1][1] = 0001000000001100000100000000101001
aPath[1][2] = 000100000001100000100000000101010
aPath[1][3] = 00010000000110000010000000101011
aPath[1][4] = 00010000000011000000100000000100100
aPath[1][5] = 00010000000110000010000000101101
aPath[1][6] = 00010000000110000010000000101110
aPath[1][7] = 00010000000011000000100000000101111
```

Figure 20. aPath after having Processed 43 Symbols

## 5 References

[1] DN504 FEC Implementation (swra113.pdf)
[2] CC430 User's Guide (slau259.pdf)
[3] CC1100 Single-Chip Low Cost Low Power RF-Transceiver, Data sheet (cc1100.pdf)
[4] CC1100E Low-Power Sub-GHz RF Transceiver (470-510 MHz \& 950-960 MHz) (CC1100E.pdf)
[5] CC1101 Single-Chip Low Cost Low Power RF-Transceiver, Data sheet (cc1101.pdf)
[6] CC1110Fx/CC1111Fx Low-Power Sub-1 GHz RF System-on-Chip (SoC) with MCU, Memory, Transceiver, and USB Controller (cc1110f32.pdf)
[7] CC1150 Single Chip Low Cost Low Power RF-Transmitter (cc1150.pdf)
[8] CC2500 Single-Chip Low Cost Low Power RF-Transceiver, Data sheet (cc2500.pdf)
[9] CC2510Fx/CC2511Fx Low-Power SoC (System-on-Chip) with MCU, Memory, 2.4 GHz RF Transceiver, and USB Controller (cc2510f32.pdf)
[10] CC2550 Low-Cost Low-Power 2.4 GHz RF Transmitter (cc2550.pdf)

## 6 General Information

6.1 Document History

| Revision | Date | Description/Changes |
| :--- | :--- | :--- |
| SWRA313 | 2010.01 .25 | Initial release. |

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[^0]:    ${ }^{1}$ The division is a "whole number" division; i.e., all variables are of type unsigned short

